# Towards Managing Variability using Software Product Family Architecture Models and Product Configurators<sup>1</sup>

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### **Abstract**

In this paper, we study the possibility of applying configurator tools developed for configuring mechanical and electronics products to representing and managing the variation points in a software product family. We compare at the conceptual level software architecture description languages and configuration modelling. Based on the analysis we are able to define a way of representing much of the architectural knowledge using the configuration modelling concepts. Thus, it seems feasible to provide software configuration support using configurators if the software is represented through architectural descriptions. However, there are also some differences that require extending the current conceptualisations of configuration knowledge and tools to capture software products adequately.

#### 1. Introduction

Software product families (or lines) have been proposed as a means for increasing the efficiency of software development and controlling complexity and variability of software products. They have become increasingly important in the software industry [1,2]. A *software product family* can be roughly defined to consist of a common architecture, a set of reusable assets used in systematically producing, i.e. deploying, products belonging to the family, and the set of products thus produced.

In this paper we pursue one approach to modelling and managing the variability of software product families, based on viewing them as configurable software product families. A configurable product is such that each product individual is adapted to the requirements of a particular customer order on the basis of a predefined configuration model [3]. Such a model explicitly and declaratively describes the set of legal product variants by defining the components out of which it can be constructed and their dependencies on each other. A specification of a product individual, i.e., a configuration, is produced based on the configuration model and particular customer requirements in a configuration task. Efficient knowledge based systems for configuration tasks, product configurators, have recently become an important application of artificial intelli-

gence techniques for companies selling products adapted to customer needs [4,5]. Stated roughly, a configurator supports the deployment process by preventing combinations of incompatible components, by making sure that all the necessary components are included, and by deducing the consequences of already made selections of components based on the dependencies in a configuration model. They are also capable of automatically generating an entire correct configuration based on requirements posed by a user. The tools are based on declarative, unambiguous modelling methods and sound inference algorithms for these. Thus, only the configuration model needs to be created to support a new product family. There are also several reports on successful use of product configurators in practice [6,7,8].

The most systematic software product families closely resemble configurable products in that they are composed of standard re-usable assets and have a predefined architecture [1]. Our approach is based on the assumption that, although customer-specific programming may be required, a significant portion of the assets can be developed and systematically modelled in advance of their deployment.

A major effort has been spent on developing architecture description languages (ADLs) that can be used for representing these re-usable assets and software architectures. Thus, product families and ADLs seem natural counterparts in the software domain for configurable products and configuration modelling languages. There are many ADLs and large differences between them [9,10].

We study the possibility of applying methods and tools developed for modelling and configuring mechanical and electronics products to configuring software. A prerequisite for coming up with a general solution to this problem is to define a mapping from the conceptualisation of software systems to a conceptualisation of configuration knowledge. Towards this end, we analyse three prominent ADLs at the conceptual level and compare them with the major concepts used for modelling configuration knowledge. Based on the analysis and comparison, we show how to represent main concepts of ADLs using the configuration modelling concepts. In addition, we identify several

<sup>&</sup>lt;sup>1</sup> This is an extended and revised version of a paper presented in the Configuration workshop of the 15<sup>th</sup> European Conference on Artificial Intelligence (ECAI 2002).

potential needs for extending the configuration modelling concepts with ADL derived concepts.

For the purposes of this paper we concentrate on three important ADLs: Acme [11,12,13], Wright [14,15] and Koala [16,17,18]. Out of these, Acme has been designed to include features of other ADLs that its designers considered central. The relevance of Acme is further promoted by the fact that one of the goals of Acme is to serve as an interchange language for other ADLs. Wright is a widely cited ADL that has a rigorous semantics and describes behavioural aspects of software. Both the use of formal methods and description of behaviour make Wright important among ADLs. Koala is not precisely an ADL, but it is in commercial use at Philips Consumer Electronics for documenting products in a product population. Therefore, being one of the few ADLs in industrial use, Koala is an important example of the practical aspects of ADLs.

As the reference point in the comparison, we employ a configuration ontology presented by Soininen et al. [19]. This ontology synthesises prior conceptualisations of configuration knowledge [6,7,8,20]. Moreover, it is very similar to another recognised configuration ontology presented by Felfernig et al. [21]. Thus, as it seems to cover most approaches to configuration modelling, it is a natural reference point for conceptual level analysis.

The remainder of this paper is organised as follows: An overview of the central concepts in product configuration and the modelling concepts in the configuration ontology is given in Section 2. An overview of software architecture and ADLs is provided in Section 3. Section 4 introduces our framework for analysing and comparing ADLs along with the most important characteristics of three ADLs. In Section 5, a comparison between the ADLs and the concepts of the configuration ontology is presented. A mapping from the most important concepts of ADLs to the concepts of the configuration ontology is given in Section 6 and potentially needed extensions of the ontology are discussed in Section 7. We discuss our findings and previous work in Section 8 and finally give our conclusions and topics for further research in Section 9.

### 2. Product configuration

In this section, we first define the fundamental concepts of product configuration. Thereafter, we introduce a configuration ontology, i.e., a conceptualisation for modeling configuration knowledge, which conceptualises these concepts.

## 2.1. Fundamental concepts

We define a *product* as an abstract specification of an entity that a company sells. A *product instance* is a product that is to be delivered to a customer or a design of a product, which is concrete enough to serve as a specification for producing it. [22]

In the domain of configurable mechanical and electronics products, a *configurable product* (or a *product family*) is defined as a product that comprises a large number of variants and serves the specific needs of the individual

customer by allowing customer-specific adaptation of the product.

Configurable products are specified through *configuration models*, which define the basic product properties and the possibilities for tailoring them. Product instances, in turn, are specified through *configurations*, which result from completing the *configuration task*. Information systems are used to support the configuration task; such information systems are called *configurators*. As mentioned in the introduction, configurators can provide a wide range of functionality, including, e.g., making deductions based on choices a user has already made, and preventing incompatible combinations of components.

During the configuration task, *user requirements* are used to constrain the set of possible configurations represented by the configuration model until there is only one correct (with respect to the configuration model) configuration that satisfies the user requirements left.

Putting pieces together, configurable products form an approach to satisfying variable customer needs. The approach involves two separate processes: first, a product (family) is defined through a configuration model; second, configurations of the product matching specific needs of a customer are formed. Product configurators can support both of the processes. What is essential is that the effort required to create a configuration of the product is moderate compared to designing a product from scratch.

### 2.2. Configuration Ontology

Next, we will give a short overview of a de facto standard configuration ontology used for modelling configurable products. For full details, the reader should refer to [19]

The configuration model consists of a set of *types*. More specifically, a configuration model consists of: a set of component types, a set of port types, a set of resource types, a set of function types - all the above-mentioned sets of types are organised in *is-a hierarchies*, and all types can be given *attributes* that present relevant information about the types. Additionally, component and function types are organised in composition hierarchy. Finally, a configuration model includes a set of constraints. A configuration is defined as a set of instances of the types. These instances are called *individuals*.

Component types represent distinguishable wholes in products.

**Example.** Figure 1 (a) depicts a simple configuration model of a computer, and will be used as a running example to illustrate the concepts of the ontology. Figure 1 (b) contains the legend of the notation used. The notation is no standard notation, but UML extended with some additional symbols. Type names are typeset using Arial, and instance names using Courier. In the figure, there are a number of component types: Home PC, Office PC, PC etc. Further, there are a number of is-a relationships between the component types: Home PC is-a PC, SW1 is-a Software etc. ■

**2.2.1. Structure.** The structural composition of component types is modelled by means of *part definitions*. Each part

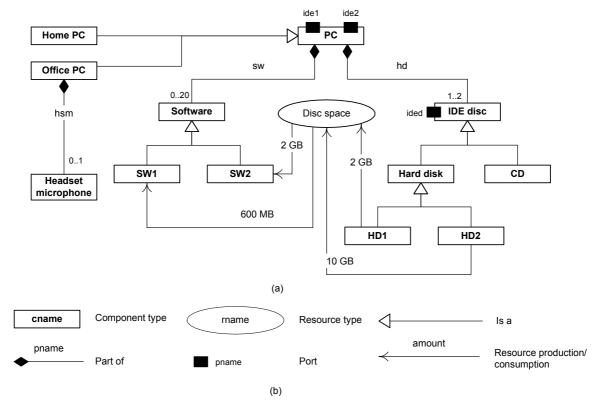


Figure 1 (a) A simple configuration model (b) Legend of the notation used

definition includes a part name, a set of possible part types, and a cardinality. The set of possible part types consists of the component types the individuals of which can occur as a part of the type containing the part definition (whole type). Further, cardinality specifies the amount of individuals that must occur as part of the whole type.

**Example.** Component types PC and Office PC include parts: the former has two parts, sw of type Software, and hd of type IDE disc; the latter has part hsm with cardinality of 0..1, which implies that the part is optional. In common terms, this means that an Office PC may, but is not required, to include a headset microphone, whereas a Home PC never includes one. ■

Alternative parts can be specified by defining multiple possible part types.

**2.2.2. Topology.** Connections between component individuals, i.e., the topology of a configuration, can be modelled by supplying component types with *port definitions*. In common terms, ports define the connection points, i.e., interfaces, components have. A port definition includes a *port name*, a *set of possible port types*, and a *cardinality*.

The semantics of a connection are that there is a physical or logical connection between the two port individuals and the two component individuals containing them.

**Example.** In Figure 1 (a), there are three ports: component type PC defines ports ide1 and ide2, and component type IDE disc defines port ided. ■

**2.2.3. Resources.** Resources are used in the ontology to model the production and use of entities, or the flow of such entities from one component individual to another.

Resource type defines the properties of a resource. A resource type defines whether the resource production or use

must be *satisfied* or *balanced*. If the production or use is to be balanced, the production must exactly match the use; in the case they must be satisfied it is enough that production is greater than or equal to use.

The production and use of resources is specified by *production definitions* and *use definitions* in component types. These definitions specify the resource types and quantities produced or used.

**Example.** There is a single resource type, namely Disc space in our sample model. This resource type is consumed by individuals of component types SW1 and SW2, and produced by individuals of HD1 and HD2. ■

**2.2.4. Functions.** All the above-presented properties of configuration models are related to technical aspects of products. However, in many cases it is necessary to communicate more abstract characterisations of the features or functionality of configurable products to salespersons and customers.

The configuration ontology includes *functions* as a means for communicating non-technical information. A *function type* is an abstract characterisation of the product. As an example, the configuration model of Figure 1 could be supplemented with the function type MusicCD with the semantics of being able to play music from CDs.

To be of any use, the function concept must be somehow related to the technical concepts. This is achieved through special constraints (see section 2.2.5 below), *implementation constraints*. They specify that a certain function be implemented by one or more technical concepts, e.g., component or port individuals. Returning to the example, MusicCD type could be implemented by individuals

of component types CD and Speaker (not in the model) together.

Similarly as for component types, function types can be defined a compositional structure through part definitions, with the natural exception that the possible part types must be function types.

**2.2.5.** Constraints. Constraints can be used to capture aspects of products that cannot be reasonably modelled using the concepts presented above. A constraint is a formal, mathematical or logical, rule specifying a condition that must hold in a correct configuration. They can be used to specify arbitrarily complex interactions of types, individuals, and their properties. Typical conditions used in constraints include *require* and *incompatible*: the semantics of the afore-mentioned are that a certain component individual in a configuration implies that another component individual must (require) or must not (incompatible) be in the configuration.

# 3. Software architectures and architecture description languages

Software architecture of a system purports to describe the high-level structure of a software system. The significance of considering architecture when designing software systems is well understood. There is, however, no single, generally accepted method for describing software architecture: UML is a good candidate for such a method, but as noted e.g. in [23], it is by no means an optimal tool for documenting all aspects of architecture. Simple methods, such as referring to an existing architectural style or using box-and-line diagrams with no or vague semantics, have been recognised to be inadequate for the task [24]. Hence, there is a need for better methods.

Architecture description languages (ADLs) are a promising candidate solution for the architecture description problem. Loosely defined, ADLs are formal notations with well-defined semantics, whose primary purpose is to represent the architecture of software systems. A large number of ADLs have been proposed. ADLs have in common the concept of component, although different ADLs have different names for the same concept [10]. But in their other characteristics, ADLs differ from each other radically. Some of them address a special application domain and others are dedicated to a specific architectural style [10]. ADLs also employ different formalisms for specifying semantics, and there is variety in how rigorously the syntax and semantics are defined.

The most fundamental elements of architectural descriptions include *components*, *connectors* and their *configurations* [10,11,24].

Components represent the main computational elements and data stores of the system. Intuitively, they correspond to the boxes in the box-and-line diagrams. Clients, servers and filters are examples of components. In a working system, a component might manifest itself as an executable file or a dynamic link library. [11]

Unlike components, connectors are not loci of application specific computation in software systems. Instead, they represent interactions between components. In a boxand-line diagram, connectors are depicted as lines between the boxes. Examples of connectors include method invocation, pipes and event broadcast. [11]

Components can be connected to each other to form configurations. They are sometimes referred to as systems [11] or architectural configurations [10]. In many ADLs, components can only be connected through connectors; explicit use of connectors has even been proposed a defining characteristic of an ADL [10]. Typically, components are connected to each other through *connection points*. Different ADLs call these connection points with different names, e.g. port, role or interface.

In some ADLs, components can also have an inner structure. Such components are called *compound components* and they represent a subsystem that has an architecture of its own. With composite components it is important to be able to specify how the inner parts of the component are linked to the component itself. Usually, the linkage is defined by binding connection points of the compound component with connection points of its parts. Intuitively, binding means that the connection point of the compound component is in fact a connection point of some other component inside the compound component.

A practical concern with ADLs is the tool support available for them. Tool support is out of the scope of this paper, since the goal is to analyse the modelling languages. However, it should be noted that support for generating executable systems out of architectural descriptions is one of the goals of research on ADLs [10]. This is a goal shared by research on configuration modelling.

# 4. Analysis of three architecture description languages

In this section, we first define a framework for analysing and comparing the concepts of ADLs with those of configuration. Thereafter, we use the framework to study three ADLs: Acme [11,12,13], Wright [14,15] and Koala [16,17,18]. A more detailed analysis can be found in [25].

#### 4.1. Framework for analysis and comparison

The fundamental phenomena described by the configuration ontology and that presented in [21] are: taxonomies, structure, topology, resources, functions, and constraints.

In the following three subsections, we will analyse the above-mentioned ADLs using a comparison framework composed of three parts. The first part includes the key concepts of ADLs and the configuration ontology, and the relations between them. The concepts include *components*, *connectors*, *configurations*, *connection points*, *attributes*, *resources*, *functions*, and *constraints*. The relations include *topology*, *taxonomy*, and *structure*. The second part considers the existence of different concepts for types and instances. The last part of the framework is the variation mechanisms provided by ADLs and the configuration ontology.

#### 4.2. Acme

The basic concepts of Acme are *components*, *connectors*, and *systems*. System is the Acme term for configuration. On the other hand, there are no constructs for resources or functions in Acme. Both components and connectors have connection points that are called *ports* for components and *roles* for connectors. *Design elements* include component, connector, port, and role. Components are connected to connectors by defining an *attachment* between the port of a component and the role of a connector. One connector may connect multiple components. Components cannot be connected directly to each other and neither can a connector to another connector. [11]

Components and connectors can have attributes that are called properties in Acme. Properties are uninterpreted values, i.e. they do not have any semantics defined.

In Acme, design constraints can be defined using first order predicate logic. They can be either *invariant* or *heuristic*: invariant constraints must hold, whereas heuristic constraints are merely hints of what should be true for an Acme system. Constraints can be used to express various aspects of Acme systems: e.g. the existence and values of properties and the connections present in a system. [12]

In addition, Acme includes a structure called *representations* that can be used for describing an alternative view of a component or a connector. *Rep-maps*, or in other words, *representation maps*, can be used to specify the correspondences between different representations of a design element. There is, however, no semantics defined for either representations or rep-maps. One possible use of these constructs is representing the compositional structure of a component and the correspondences between the ports and roles of the compound component and those of the contained components. [11].

Although types are not first class entities in Acme, it has two type systems: one for design element types and, and another for systems. Types in the design element type system are sets of required structure, i.e. design element declarations, and values. New types can be formed from existing types through subtyping. System types are called *families*. A family consists of design element type definitions. Subtypes of families can be formed through single or multiple inheritance. Also, a system can be declared to be a member of many family types. [13]

What makes types a secondary concept in Acme is that design elements and systems need not have a type or be a member of a family, respectively. A design element being of a given type merely implies that the design element has the structure and values specified by that type. Similarly with families, a system being a member of a family signals that the type definitions of the family are type definitions of the system, too. Therefore, type systems of Acme can be considered a sort of macro expansion mechanism.

The syntax and semantics of Acme are formally defined, the latter in terms of a mapping to first order predicate logic.

There seem to be no constructs in Acme for modelling variety. What seems to come closest to modelling variability is the family construct. It can be used to specify a set of

type definitions shared by a set of systems. Furthermore, constraints can be used to enforce the instantiation of certain design elements. Hence, the family definitions complemented with constraints seem to provide a mechanism for specifying product families with certain properties.

### 4.3. Wright

As in Acme, there are *components*, *connectors*, *systems*, *ports*, *roles* and *attachments* in Wright and their semantics are the same in both languages. There are no attributes, resources or functions. What distinguishes Wright from Acme and makes it special among ADLs is its way of specifying the behaviour of ports, roles, connectors and components, and the possibilities for analysis based on these specifications. Wright uses CSP (Communication Sequential Processes) [26], a formal approach for two purposes: (1) specifying processes that reside in Wright elements and (2) defining semantics of non-CSP parts of the language. CSP is a formal method for specifying and analysing the behaviour of objects in terms of sequences of events in which they engage. The pattern of events that is possible for an object is termed a *process*. [14,15]

Each port and role is associated with a CSP process. In addition, each connector and component includes a separate glue and computation process, respectively. The glue of a connector defines the operation of the connector as an entity. That is, the glue coordinates the operations of the other processes in the connector. Ports are attached to roles to form systems. Which ports can be attached to which roles, is determined by their process descriptions: a port can be attached to a role if the port will behave well in all situations enabled by the role, i.e., CSP defines a compatibility relation between ports and roles.

The second usage of CSP in Wright, defining semantics of non-CSP parts of the language, allows using tools operating on CSP to reason about properties, most notably about dead-lock freedom, of a Wright connector. This is an important class of tool support enabled by the rigorously defined semantics of Wright.

Wright allows describing hierarchical structure of both components and connectors. This is done by enclosing a system into the place of a process. In addition to the normal system specification, bindings between the port and role names in the enclosing element and those specified in the enclosed system need to be specified.

Wright distinguishes between component and connector types and instances. Each connector and component is of exactly one type. There is, however, no taxonomy of types.

In addition to component and connector types, Wright includes a construct called *style*. Styles are collections of type definitions and *constraints*. They are expressed in first order predicate logic and they can be used in a manner similar to that in Acme described above. In addition to component and connector type definitions, a style can include *interface type* definitions. They are process descriptions that can be used in port and role definitions.

Type definitions in styles can be parameterised. That is, parts of the type definition can be left open and a value can be filled in when the type is instantiated. New styles can

be defined in terms of existing ones through subtyping: the new style has the same type definitions and constraints as the old one plus some additional type definition or constraints.

Variation mechanisms of Wright are similarly limited as for Acme, although Acme uses the term family where Wright uses style. In short, styles supplemented with constraints seem to be able to express variability.

#### 4.4. Koala

As the languages described above, the Koala model has *components* as a main design element. But in other respects, Koala differs greatly from its peers. In Koala, there is no notion of connectors, resources, functions or constraints. Components are connected, or *bound*, as it is said in Koala, to each other through *interfaces* that are the connection points in Koala. The connection between components is not symmetric: a distinction is made between *provided* and *required interfaces*. Loosely defined, a component having a provided interface means that the component offers some service for other components to use. Similarly, a required interface signals a service being required by the component from some other component. Koala interfaces are similar to those in COM or Java.

Compound components can be used to express compositional structure in Koala, i.e. other components can be contained within a component. An interface of a compound component can be bound to an interface defined by a contained component. A *configuration* is defined as a component that is not contained in another component and defines no interfaces.

In addition to binding interfaces to each other, it is possible to bind constituent parts of interfaces directly. These parts are called *functions*. Hence, interfaces in Koala are not atomic even when considered as connection points.

There are some limitations on how interfaces can be connected to each other: These limitations are best illustrated with the aid of Figure 2. In the figure, components are depicted as boxes and interfaces as squares containing triangles inside them. A required interface is depicted as a triangle pointing outwards from the component, and a provided interface with a triangle pointing inwards. The binding rules are that must be bound by its tip to exactly one base of an interface, and any number of interfaces can be bound to the base of an interface. Notice that these rules cover both bindings between interfaces in contained components and interfaces in independent component, i.e., components such that neither is contained within another.

Another rule concerning the bindings is that the type of the tip interface must be a *supertype* of the type of the base interface: interface type A is a supertype of B exactly when B contains all the functions of A.

Koala has a type system: a distinction is made between both interface and component types and instances. There is, however, no taxonomy of component or interface types beyond the supertype relation mentioned above.

Koala includes a construct, *module*, which is a component without an interface of its own. Modules are used inside compound components for gluing interfaces. Sup-

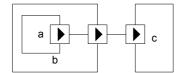


Figure 2 Koala components and interfaces

pose, for example, that each component contained in a compound component has an initialisation interface to be called before using the component. Due to binding rules, it would not be possible to bind all these interfaces to any single interface of the compound component. Therefore, a new configuration specific module is added: when the initialisation function for the compound component is called, the call is routed to the module, which in turn calls the initialisation functions of all necessary components in the order desired.

In addition to the constructs already mentioned, Koala provides mechanisms for handling both the *internal diversity* of components and the *structural diversity* in a configuration. Internal variety is manifested as variation of component parameters. There may be dependencies between parameters: a parameter value may imply that another parameter has a certain value. Structural diversity pertains to alternative provided interfaces for a required interface: e.g. there may be multiple components that provide the same interface required by a certain component. The choice between the interfaces is made by a construct called *switch* either statically, that is at compile time, if the information required for the selection is available, or, otherwise, dynamically at runtime.

# 5. Comparison of concepts of the adls with the configuration ontology

In this section, we use our framework defined in the previous section for comparing the concepts and constructs found in the ADLs with those of the configuration ontology.

### 5.1. Key concepts and relations between them

Component is the central concept of Acme, Wright and Koala. It is also present in the configuration ontology with that same name. The semantics are as well similar: components represent the defining parts of a system in configuration modelling, too. In addition, systems as defined in Acme and Wright and configurations as defined in Koala have a counterpart in the configuration ontology, namely configuration.

The notion of connection points is also common to all the studied modelling methods. In Acme and Wright they are called ports and roles in components and connectors, respectively. In Koala connection points are termed interfaces and in the ontology ports. The semantics of connection points are also similar in all the disciplines: they denote the mechanism for connecting other entities.

Connectors are first-class citizens in Acme and Wright., but there are no connectors in the configuration ontology. In Koala, modules can be considered as a form of connectors. Thus, there is a major difference in how the disciplines handle architectural connection – an important issue in both ADLs and in the configuration ontology.

What then is the reason for this disagreement in architectural connection? We believe that at least a partial reason for the importance of connectors in Acme and Wright can be found in the underlying assumptions of them and several of the ADLs not studied in this paper: a major issue in software architecture has been reusing existing components. Furthermore, there has been considerable effort in the software engineering community to reuse heterogeneous components, which cannot be connected directly to each other due to different communication mechanisms and various other reasons. Therefore, connectors have been introduced in ADLs as a vehicle for connecting heterogeneous components.

In Koala, the situation is rather different: components are generally rather homogenous, and it seems that there is no need to require that connectors be used whenever connecting components: modules are used as needed. Hence, in this respect, Koala is closer to the configuration ontology than Wright and Koala.

Resources, a feature present in the configuration ontology but not in any of the ADLs, is similar to the notion of provided and required interfaces present in Koala in the sense that they are both anti-symmetric. What is more, resources are produced and consumed by components, just as interfaces are provided and required. However, resources are produced and consumed in certain quantities, which gives them more expressive power compared with the notion of provided and required interfaces.

In addition to simulating provided and required interfaces, resources can be used to model other relevant quantities. Such quantities include memory, power, output capacity and throughput. The software engineering community has considered similar issues important [27]. Hence, resources could very well be an important feature of the configuration ontology when used to model software architecture.

Modelling functions is another feature of the configuration ontology that all three ADLs presented in this paper lack. Functions are an important aspect of software engineering usually termed features in the domain [28]. We believe that also functions could be very useful when modelling software with the ontology.

All the ADLs have some mechanisms for modelling structure. However, the configuration ontology provides much stronger mechanisms: the configuration ontology provides a wide range of variation mechanisms. Furthermore, in the configuration ontology a component can be a part of many components simultaneously, which is not possible in any of the ADLs.

All the disciplines except Koala have explicit mechanisms for expressing constraints. Further, in all disciplines where constraints exist, they are logical expressions about the non-behavioural properties of a system modelled in that discipline. A difference is that in the configuration ontology, there is no direct support for heuristic constraints as defined in Acme. Support for modelling preferences and optimisation criteria have been identified as important and developed in other research on configuration.

## 5.2. Distinction between types and instances

All the three ADLs have some distinction between types and instances. In Acme, the distinction is rather weak, as the type systems can be seen as a simple macro expansion mechanism. Nevertheless, there is taxonomy between the Acme types. The situation is rather similar in Wright: types bear a little meaning as such. The only function of types seems to be facilitating in defining and altering recurring patterns. In Koala, the supertype relation constraints bindings between interfaces. This relation is not, however, declared explicitly, but implicitly based on interface types. The component types seem to have no function beyond defining the structure of a set of components. Hence, component types seem to be as a construct as weak as types in Acme and Wright.

In the configuration ontology, strong distinction between types and instances is one of the basic assumptions and is made for all kinds of entities. Types are organised in taxonomies.

#### 5.3. Variation mechanisms

A question closely related with the distinction of types and instances is: What is being modelled, one product or a product family. The configuration ontology aims at modelling product families. Configuration model knowledge defines the common properties of the family members. A lot of variation mechanisms are provided.

As stated in the analysis of Acme and Wright, both of these languages can be seen to provide some support for modelling variability: there are no explicit variation mechanisms, but the combination of system types and constraints seem to be able to express common structure shared by a set of products.

In Koala, there is some knowledge about the common properties of all the products: component and interfaces definitions are stored in a component repository and they are common to different systems to be constructed [17]. In fact, type definitions shared by a set of products is exactly the same phenomenon be have already seen in family construct Acme and in the style construct in Wright. As there are no constraints to complement the shared type definition in Koala, the support provided by Koala for variability is weaker than that Acme and Wright.

In the previous section, it was stated that Koala could model both internal and structural variety. How does this statement relate to the above observation that Koala provides a weaker support for variability than Acme and Wright? We claim that we are dealing with two distinct forms of variability. The variability in Acme and Wright can be used to span a set of products with many similarities, or a product family. On the other hand, the variation mechanisms in Koala seem to model behavioural variety of software embedded in a physical product instance: e.g. a television set can behave differently depending on some parameters. Of course, it could be argued that the television set in our example is, in fact, a product family. However, we consider the variation mechanisms discussed above examples of different phenomena.

# 6. Modelling software architecture with the configuration ontology

In this section, we strive to synthesise the configuration ontology with the domain of software architecture. We do this by mapping the concepts in the ADLs to some concept or concepts in the configuration ontology. Components, ports, properties, and constraints are represented in the obvious manner using their direct counterparts, whereas the representation of connectors and roles is more problematic. Hence, we will present a mapping of connectors to components, and provided and required interfaces to ports with the aid of type specifications.

Due to limited space, we will only present the main ideas of the mapping. For full details, please refer to [25].

# 6.1. Modelling connectors as a type of component

In translating the semantics of connectors in Acme and Wright into concepts in the configuration ontology, it helps to observe that components and connectors have structures very similar to each other. Therefore, it is natural to view connector as a subtype of component with special semantic constraints. Indeed, defining connector to be a subtype of component will enable us to express part of the semantics associated with connectors. Furthermore, we can define roles in connectors to be ports in the connector-type components. To enforce the right use of connectors, we define suitable constraints that enforce the right use of connectors: e.g. in Wright, the only class of allowed connections is that between a component and a connector.

Subtyping can also be used for distinguishing provided and required interfaces from one other. By defining common supertypes for provided and required interfaces it is possible through multiple inheritance to have two versions of each port type, a provided and a required. By using constraints it is possible to assert that invariants concerning provided and required interface types hold. For instance, the fact that in Koala a required interface must be connected to exactly one provided interface of the same interface type can be easily captured using constraints.

### 6.2. Capturing diversity

Internal diversity of Koala can be captured with attributes defined by components and constraints. Dependencies between different parameters can be captured using constraints between attribute values of component types.

In the configuration ontology, cardinality of a port defines the amount of ports that can be connected to it. Cardinality can be used to capture some aspects of structural diversity in Koala. By defining cardinality greater than one for a port representing a required interface, multiple provided interfaces represented as ports could be connected to that port. This is only a partial solution as it says nothing about deciding which ports should actually be connected; constraints can be used to model this.

# 7. Extensions needed for modelling software architecture

Albeit the ontology captures a major part of aspects of all the studied ADLs, each of them has some features the modelling of which would require extending the ontology.

Capturing all of the idea behind **heuristic constraints** of Acme may require adding some method of representing optimisation criteria and preferences in the ontology.

There is no mechanism in the configuration ontology for **modelling behaviour** similar to the way how CSP is used in Wright. In fact, the configuration ontology ignores behavioural aspects entirely. In case considering behaviour should be required in the configuration ontology, it would be natural to extend the constraint language to cover behavioural aspects, as the constraint language can be seen as the extension mechanism of the ontology.

Koala includes the method of **function binding**, in which the constituent functions of interfaces are connected directly to each other instead of connecting interfaces [17]. This construct gives an internal structure to Koala interfaces. Given that interfaces of Koala are modelled with ports in the configuration ontology, this contradicts with the underlying assumption of ports being undividable connection points. As a result, there is a mismatch between interfaces in Koala and ports in the configuration ontology.

There is a number of possible ways to capture ports with internal structure. The first one is to make Koala functions the basic level of connection. Unfortunately, this approach introduces major problems. Firstly, interfaces would lose their counterpart in the configuration ontology. Secondly, following the approach would likely lead to increased complexity in models of software products: that an interface can contain several functions implies this.

The second approach would be to introduce compositional structure for ports of the configuration ontology. Applied to the problem at hand, interface types correspond to port types that have ports corresponding to functions as their parts. This approach is appealing: it models the relation between interfaces and functions in a way corresponding to the intuitive understanding of the issue. This approach would require major changes to the ontology.

Binding of interfaces of a compound component with the interfaces of the inner parts is another feature of Koala lacking a counterpart in the ontology. It seems that the ontology would need to be extended in order for it to model this phenomenon.

## 8. Discussion, comparison to previous work

There is an apparent difference in the natures of the sets of product variations modelled in different disciplines. In the configuration domain, this set is typically termed as configurable product or a product family. One of the defining characteristics of this concept is a pre-designed general structure with a lot of variation in the configurations [29]. On the other hand, it was found that Koala supports no common structure for a set of products. In fact, Koala is not targeted at modelling a product family or a set of them, but product populations, defined as a sets of products with

many commonalities but also with many differences [16]. Hence, the underlying aims of the ADL modelling and configuration modelling are not totally similar.

In the previous section, it was stated that no satisfactory mapping could be found for function binding of Koala. One possibility to respond to this and similar problems is to ignore the problematic feature. Even though we do not light-heartedly ignore aspects of ADL that are of practical or theoretical importance, we still believe that doing so in some cases will increase the usefulness of the configuration ontology in modelling software products. Therefore, the question is: which features of ADLs should be modelled. This question can only be fully answered by empirically studying software product families.

Research closely related to this paper has been conducted earlier. We do not know of earlier attempts of comparing the concepts of software architecture description to those of configuration modelling. This is the main contribution of our paper: studying the feasibility of configuration techniques to software variability management.

In their work, Männistö et al. have pointed out the existence of the research area of configurable software and identified some key concerns in the area [30]. They have not, however, studied the concepts of ADLs in detail or proposed any mapping from these concepts to those of configuration modelling domain.

On the other hand, [31] presents a formalised software configuration management (SCM) ontology. The concepts of the SCM ontology are, however, different from those of the configuration ontology. They are aimed at representing the modules, files, or packages, their versions and the dependencies between these. The ontology does not take into account the connections and interfaces between components of a system.

Felfernig et al. have proposed a scheme for constructing configurators based on UML descriptions of configuration knowledge [21]. Their approach could be used for creating configurators for software products as well. Their approach is, however, different from our approach: theirs is based on presenting configuration knowledge in UML, while our approach is based on modelling software with the concepts of product configuration.

In [32], Kühn has presented an approach to software configuration based on structure and behaviour. He uses statecharts, a method similar to finite state machines, for specifying the behaviour of a module. This approach is similar to Wright in that it describes both structure and behaviour. With its focus on using behavioural constraints for making decisions during the configuration process, this approach is different from ours.

Feature models have been suggested as a modelling method for software product lines (see, e.g., [33]). Apparently, feature models share much with the configuration modelling concepts presented in this paper: features, and both components and functions in the configuration ontology are organised in composition hierarchies. What seems to be different in the two approaches is that the configuration ontology distinguishes between technical and nontechnical aspects of a product, whereas in features models, both aspects are contained in the same hierarchy.

Lars Geyer et al. have identified the need and enumerated a number of requirements for a configuration technique for software product families [34]. Of the requirements, the configuration ontology supports hierarchical structures for both technical and non-technical entities, i.e., components and functions, respectively. Further, the ontology incorporates a constraint mechanism. Finally, a prototype tool for configuring mechanical products that supports a subset of the configuration concepts presented in this paper corresponding to the requirements posed by Geyer et al., visualises the aggregation hierarchy and is able to assist in configuring the product, as required by them [35]. Therefore, although Geyer et al. have deemed knowledge-based configuration techniques rather useless in the context of software product families, we feel that knowledge-based configuration techniques provide considerable potential in domain of software product families.

#### 9. Conclusions and further work

We have presented an analysis and comparison of three ADLs and a conceptualisation of configuration knowledge. We defined a mapping from the concepts of ADLs to those of the conceptualisation of configuration knowledge. Our goal is to use the configuration concepts and their supporting tools for configuring software product families.

We found counterparts and close correspondences in the configuration ontology for the main elements of the ADLs we have studied and were able to propose a mapping between them that shows that configuration languages can be used for representing architectural knowledge. For instance, both share the notion of components. Furthermore, compositional structure, systems formed of connected components and constraints are phenomena present in both disciplines. Hence, it seems that the concepts of the configuration ontology can be used for modelling software products.

Capturing some aspects of ADLs seems to require extending the configuration ontology. These aspects include function binding and binding the connection points of compound components with connection points in its inner parts. Another important aspect is modelling behaviour. Of the ADLs, Wright models behaviour. Additionally, the approach presented by Kühn also emphasises behaviour [32]. The question whether behavioural aspects really are important and should be modelled when configuring software product families, should be resolved through empiric studies with real products. The existence of Koala, a commercial ADL with no behaviour modelling, suggests that modelling behaviour is not absolutely necessary.

There are still open questions and a need for further work. An ontology and a configuration language for software products should be defined, and a configurator supporting this language constructed. This work is currently in progress. This will probably require investigating more thoroughly the current ADLs and the conceptualisations of disciplines such as SCM, generative and feature based programming [28,36], and, of course, the developments in the UML community, as well as case studies of real software product families. After completing this, case studies

are needed to verify the applicability of the configuration language to modelling software. Another issue to be concerned is the computational complexity of configuring software products. Theoretical complexity analysis can provide insight into this issue, but only experiments with real products will give relevant information on the practical feasibility from this point of view. When moving towards empirical studies, it is also necessary to consider which of the existing configurators and their modelling languages best support software configuration at a more detailed level than in this study.

Finally, the economics of our approach should be studied. Particularly for simpler products, the overhead from developing a configuration model of a software product family using a configurator is higher than the advantages that can be gained. However, we believe that there are cases where the product family is complex enough, including even thousands of variation points, that the support for deployment process would outweigh the costs.

### Acknowledgements

We gratefully acknowledge the financial support from Academy of Finland (project 51394) and National Technology Agency of Finland (Tekes).

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